

# Mealy Automata, Orbit Trees and the Burnside Problem

Thibault Godin  
Highlights 2016 Brussels



ANR JCJC 12 JS02 012 01

# The Burnside problem



## Burnside Problem (1902):

Can a finitely generated group have all elements of finite order and be infinite?

# The Burnside problem

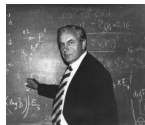


## Burnside Problem (1902):

Can a finitely generated group have all elements of finite order and be infinite?

$$G = \langle a_1, a_2, \dots, a_n \rangle = \langle \Sigma \rangle ; g = w \in \Sigma^* ; \\ \forall g, \exists n, g^n = e$$

# The Burnside problem



## Burnside Problem (1902):

Can a finitely generated group have all elements of finite order and be infinite?

$$G = \langle a_1, a_2, \dots, a_n \rangle = \langle \Sigma \rangle ; g = w \in \Sigma^* ; \\ \forall g, \exists n, g^n = e$$

Golod and Shafarevich: Yes! (1964)

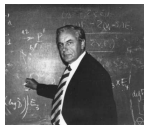
# The Burnside problem



## Burnside Problem (1902):

Can a finitely generated group have all elements of finite order and be infinite?

$$G = \langle a_1, a_2, \dots, a_n \rangle = \langle \Sigma \rangle ; g = w \in \Sigma^* ; \\ \forall g, \exists n, g^n = e$$



Golod and Shafarevich: Yes! (1964)



Grigorchuk: example generated by a Mealy automaton (1980)

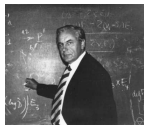
# The Burnside problem



## Burnside Problem (1902):

Can a finitely generated group have all elements of finite order and be infinite?

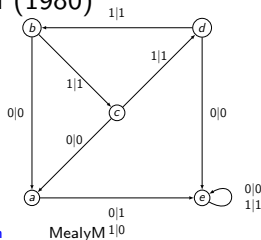
$$G = \langle a_1, a_2, \dots, a_n \rangle = \langle \Sigma \rangle ; g = w \in \Sigma^* ; \\ \forall g, \exists n, g^n = e$$



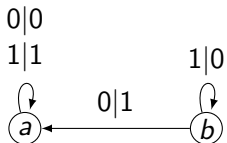
Golod and Shafarevich: Yes! (1964)



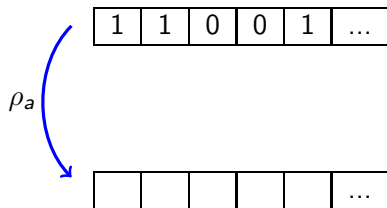
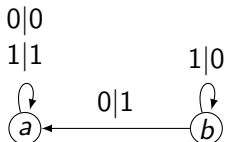
Grigorchuk: example generated by a Mealy automaton (1980)



# How to Generate Automaton Groups

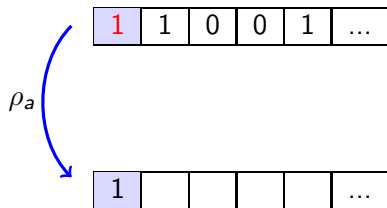
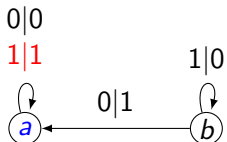


# How to Generate Automaton Groups

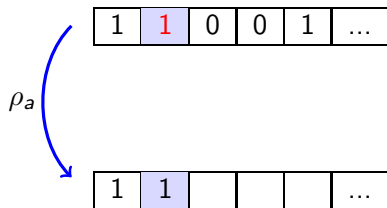
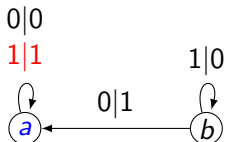




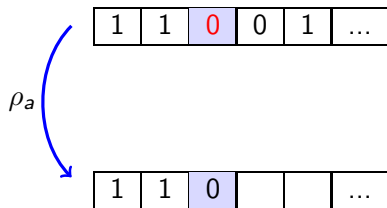
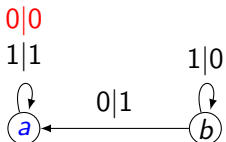
# How to Generate Automaton Groups



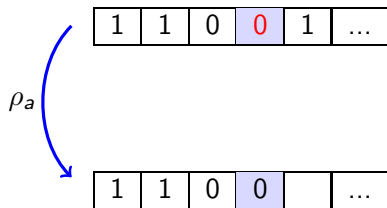
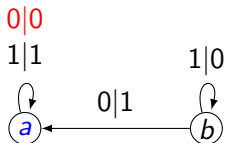
# How to Generate Automaton Groups



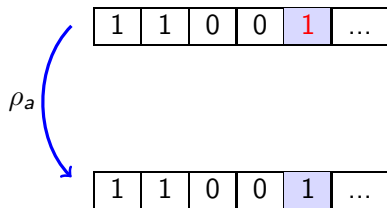
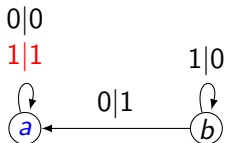
# How to Generate Automaton Groups



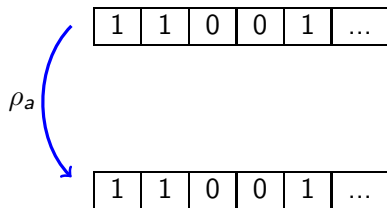
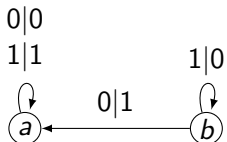
# How to Generate Automaton Groups



# How to Generate Automaton Groups

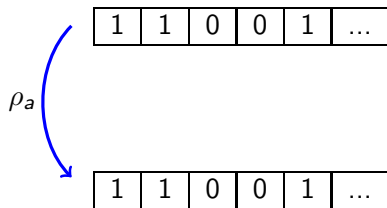
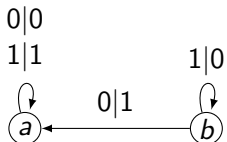


# How to Generate Automaton Groups



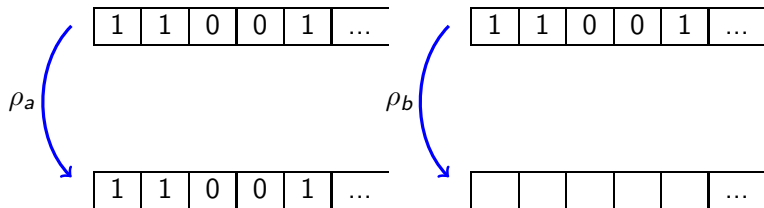
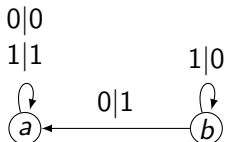
$$\rho_a : \Sigma^* \rightarrow \Sigma^*$$

# How to Generate Automaton Groups



$$\rho_a : \Sigma^* \rightarrow \Sigma^*$$
$$u \mapsto u$$

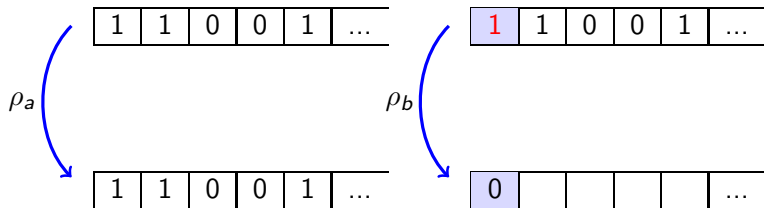
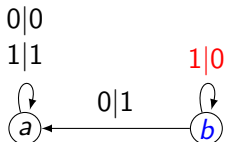
# How to Generate Automaton Groups



$$\rho_a : \Sigma^* \rightarrow \Sigma^*$$
$$u \mapsto u$$

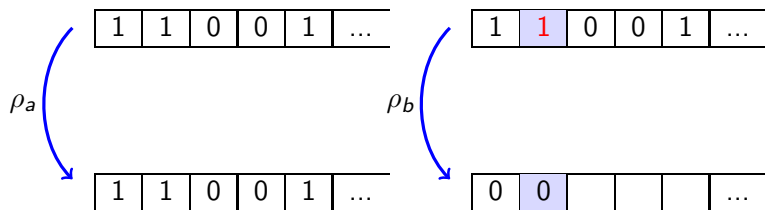
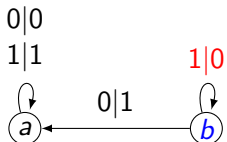


# How to Generate Automaton Groups



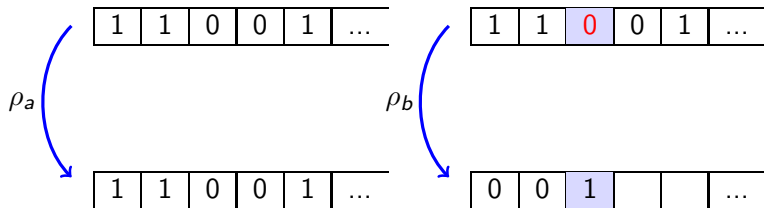
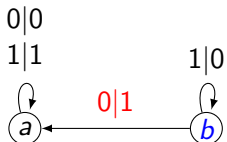
$$\rho_a : \Sigma^* \rightarrow \Sigma^*$$
$$u \mapsto u$$

# How to Generate Automaton Groups



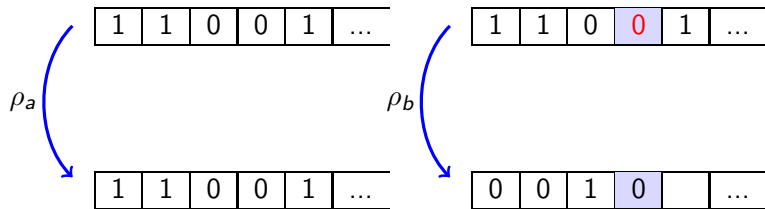
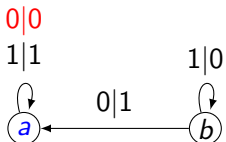
$$\rho_a : \Sigma^* \rightarrow \Sigma^*$$
$$u \mapsto u$$

# How to Generate Automaton Groups



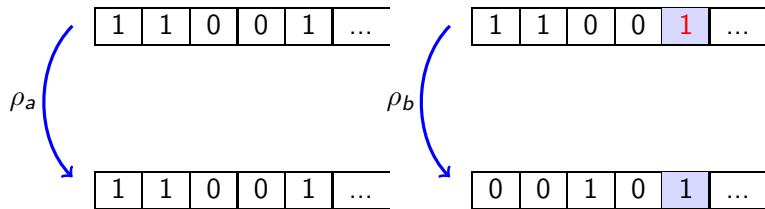
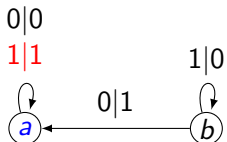
$$\rho_a : \Sigma^* \rightarrow \Sigma^*$$
$$u \mapsto u$$

# How to Generate Automaton Groups



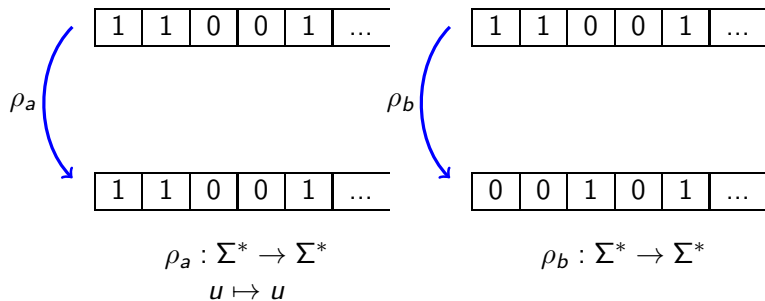
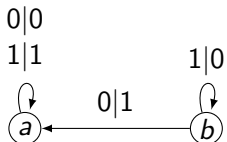
$$\rho_a : \Sigma^* \rightarrow \Sigma^*$$
$$u \mapsto u$$

# How to Generate Automaton Groups

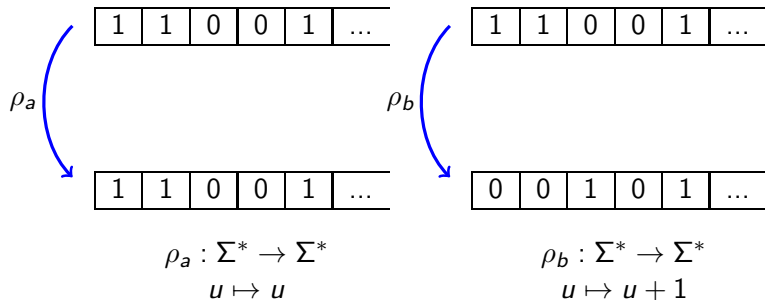
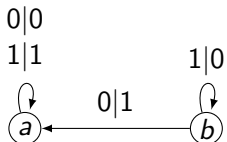


$$\rho_a : \Sigma^* \rightarrow \Sigma^*$$
$$u \mapsto u$$

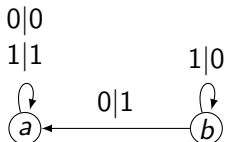
# How to Generate Automaton Groups



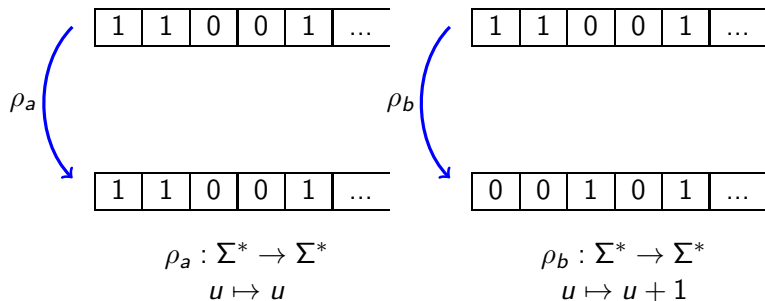
# How to Generate Automaton Groups



# How to Generate Automaton Groups

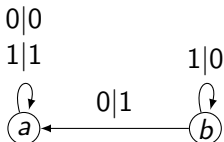


$$\langle \mathcal{A} \rangle = \langle \rho_a, \rho_b \rangle = \langle \rho_a, \rho_b, \rho_a^{-1}, \rho_b^{-1} \rangle$$

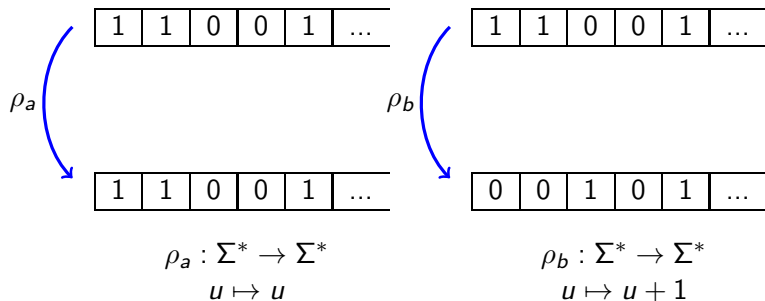




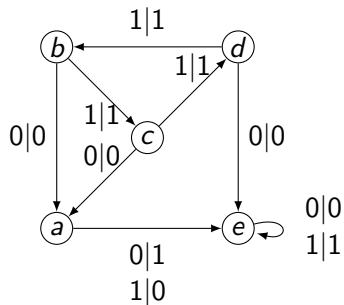
# How to Generate Automaton Groups



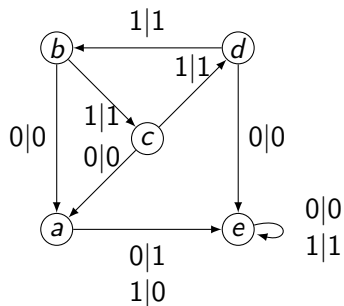
$$\langle \mathcal{A} \rangle = \langle \rho_a, \rho_b \rangle = \langle \rho_a, \rho_b, \rho_a^{-1}, \rho_b^{-1} \rangle \simeq \mathbb{Z}$$



## About the Grigorchuk automaton



# About the Grigorchuk automaton



Action of a state on the letters:

$$\rho_a : 0 \mapsto 1 \mapsto 0$$

$$\rho_b : 0 \mapsto 0; \quad 1 \mapsto 1$$

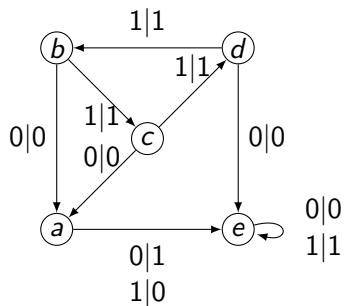
$$\rho_c : 0 \mapsto 0; \quad 1 \mapsto 1$$

$$\rho_d : 0 \mapsto 0; \quad 1 \mapsto 1$$

$$\rho_e : 0 \mapsto 0; \quad 1 \mapsto 1$$

→ permutations

# About the Grigorchuk automaton



Action of a state on the letters:

$$\rho_a : 0 \mapsto 1 \mapsto 0$$

$$\rho_b : 0 \mapsto 0; \quad 1 \mapsto 1$$

$$\rho_c : 0 \mapsto 0; \quad 1 \mapsto 1$$

$$\rho_d : 0 \mapsto 0; \quad 1 \mapsto 1$$

$$\rho_e : 0 \mapsto 0; \quad 1 \mapsto 1$$

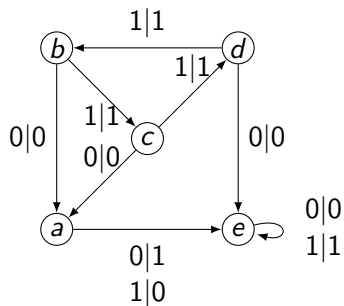
→ permutations

Action of a letter on the states:

$$\delta_0 : a, d, e \mapsto e; \quad b, c \mapsto a$$

→ not a permutation

# About the Grigorchuk automaton



Action of a state on the letters:

$$\rho_a : 0 \mapsto 1 \mapsto 0$$

$$\rho_b : 0 \mapsto 0; \quad 1 \mapsto 1$$

$$\rho_c : 0 \mapsto 0; \quad 1 \mapsto 1$$

$$\rho_d : 0 \mapsto 0; \quad 1 \mapsto 1$$

$$\rho_e : 0 \mapsto 0; \quad 1 \mapsto 1$$

→ permutations

→ **invertible**

Action of a letter on the states:

$$\delta_0 : a, d, e \mapsto e; \quad b, c \mapsto a$$

→ **not a permutation**

→ **non-reversible**

## Theorem 1 [Klimann STACS '13]

A 2-state invertible-reversible Mealy automaton cannot generate an infinite Burnside group

### Theorem 1 [Klimann STACS '13]

A 2-state invertible-reversible Mealy automaton cannot generate an infinite Burnside group

### Theorem 2 [Klimann, Picantin and Savchuk DLT' 15]

A connected 3-state invertible-reversible Mealy automaton cannot generate an infinite Burnside group





## The Burnside problem and Reversible automata

### Theorem 1 [Klimann STACS '13]

A 2-state invertible-reversible Mealy automaton cannot generate an infinite Burnside group

### Theorem 2 [Klimann, Picantin and Savchuk DLT' 15]

A connected 3-state invertible-reversible Mealy automaton cannot generate an infinite Burnside group

### Theorem 3 [G., Klimann and Picantin, LATA '15]

An invertible-reversible Mealy automaton with a non coreversible component cannot generate an infinite Burnside group

### Theorem 1 [Klimann STACS '13]

A 2-state invertible-reversible Mealy automaton cannot generate an infinite Burnside group

### Theorem 2 [Klimann, Picantin and Savchuk DLT' 15]

A connected 3-state invertible-reversible Mealy automaton cannot generate an infinite Burnside group

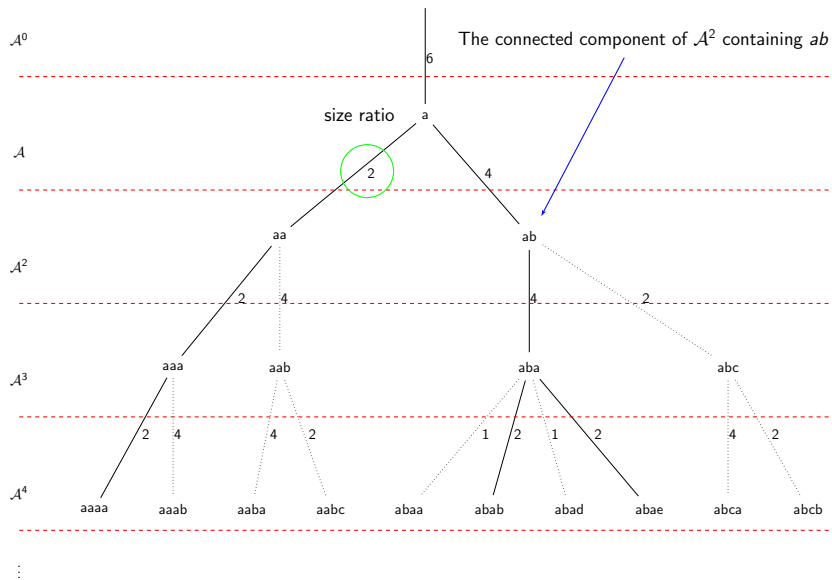
### Theorem 3 [G., Klimann and Picantin, LATA '15]

An invertible-reversible Mealy automaton with a non coreversible component cannot generate an infinite Burnside group

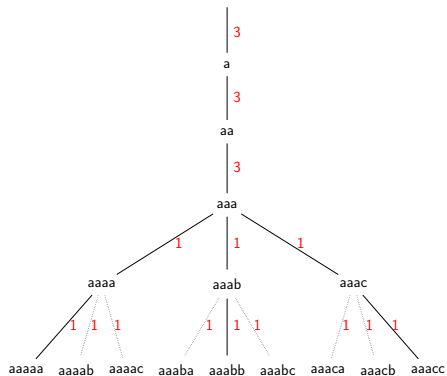
### Theorem 4 [G., Klimann MFCS'16]

A connected invertible-reversible Mealy automaton with a prime number of states cannot generate an infinite Burnside group.

# Orbit Tree



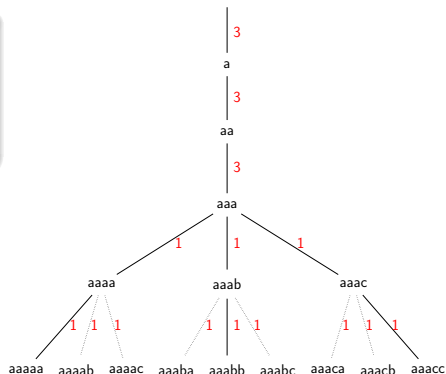
# Boundedness



# Boundedness

## Prop

$\langle \mathcal{A} \rangle$  is finite iff the sizes of the connected components of  $(\mathcal{A}^n)_n$  are bounded



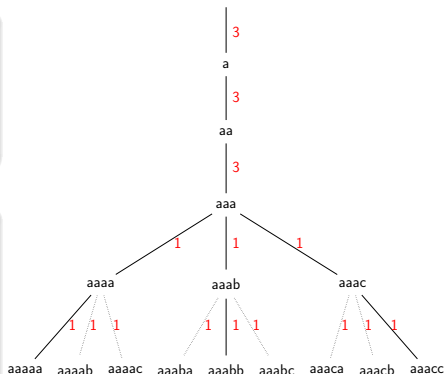
# Boundedness

## Prop

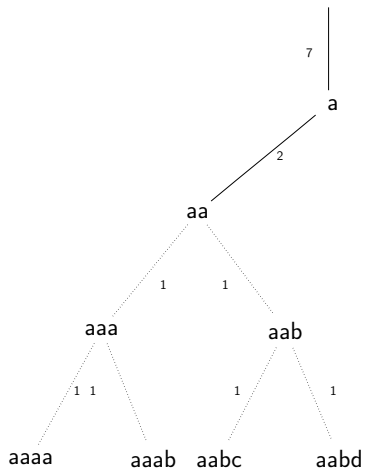
$\langle \mathcal{A} \rangle$  is finite iff the sizes of the connected components of  $(\mathcal{A}^n)_n$  are bounded

## Prop

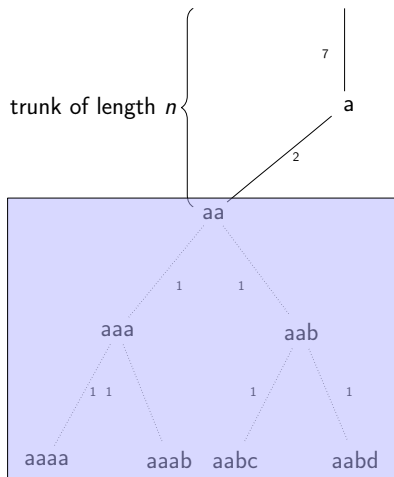
$\rho_u$  has finite order iff the sizes of the connected components of  $(\mathcal{A}^n)_n$  containing  $u^n$  are bounded



# Jungle tree

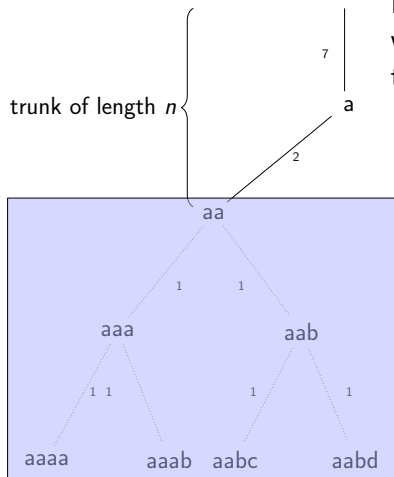


# Jungle tree



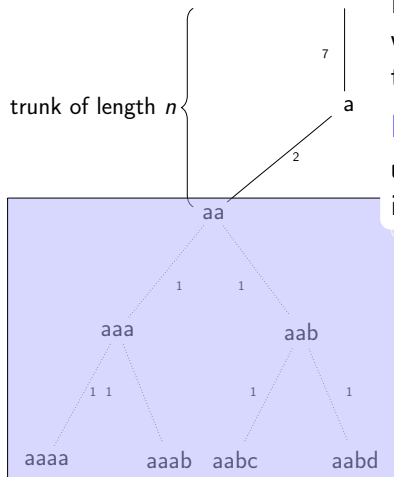


# Jungle tree



Idea: for any word, find a word with the same action in the jungle-tree

# Jungle tree

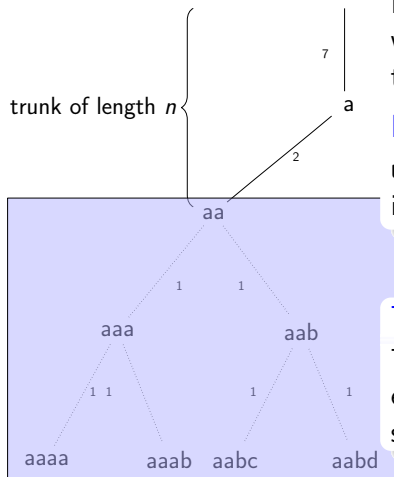


Idea: for any word, find a word with the same action in the jungle-tree

## Lemma

$\mathbf{u} \sim \mathbf{v} \Leftrightarrow \mathbf{u}\mathbf{w}\mathbf{v} \in \mathcal{J}$  and  $\rho_{\mathbf{u}\mathbf{w}} = id$   
is an equivalence relation on  $Q^n$

# Jungle tree



Idea: for any word, find a word with the same action in the jungle-tree

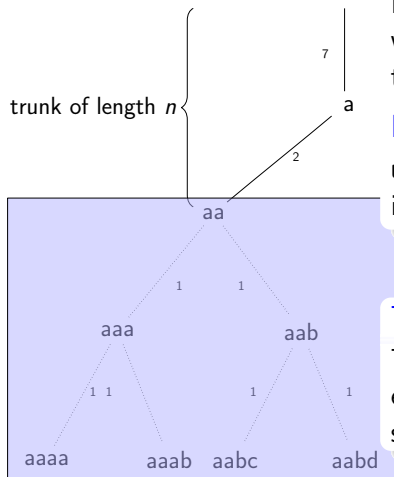
## Lemma

$\mathbf{u} \sim \mathbf{v} \Leftrightarrow \mathbf{u}\mathbf{w}\mathbf{v} \in \mathcal{J}$  and  $\rho_{\mathbf{u}\mathbf{w}} = id$   
is an equivalence relation on  $Q^n$

## Theorem

The set of first letter of an equivalence class is the whole stateset

# Jungle tree



Idea: for any word, find a word with the same action in the jungle-tree

## Lemma

$u \sim v \Leftrightarrow uvv \in \mathcal{J}$  and  $\rho_{uv} = id$  is an equivalence relation on  $Q^n$

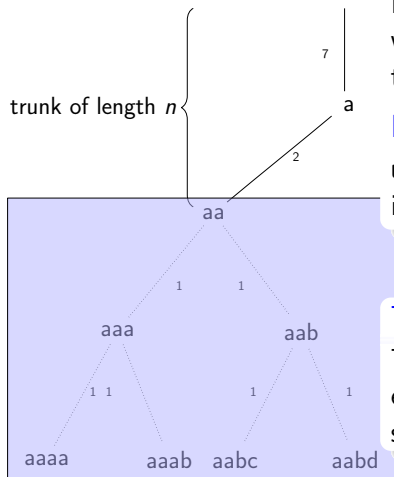
## Theorem

The set of first letter of an equivalence class is the whole stateset

$$u_1 u_2 \dots \equiv v w u_1$$

with  $\rho_{vw} = id$

# Jungle tree



Idea: for any word, find a word with the same action in the jungle-tree

## Lemma

$u \sim v \Leftrightarrow uvv \in \mathcal{J}$  and  $\rho_{uv} = id$  is an equivalence relation on  $Q^n$

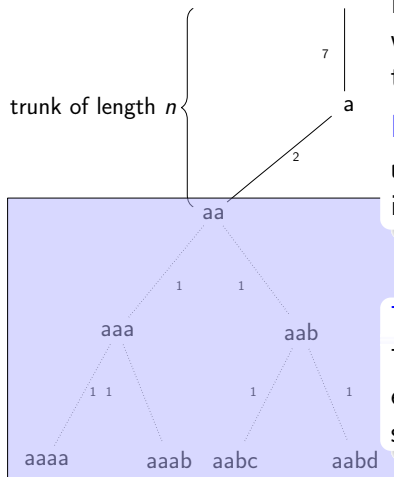
## Theorem

The set of first letter of an equivalence class is the whole stateset

$$u_1 u_2 \dots \equiv v w u_1 v'$$

with  $\rho_{vw} = id$

# Jungle tree



Idea: for any word, find a word with the same action in the jungle-tree

## Lemma

$\mathbf{u} \sim \mathbf{v} \Leftrightarrow \mathbf{u}\mathbf{w}\mathbf{v} \in \mathcal{J}$  and  $\rho_{\mathbf{u}\mathbf{w}} = id$  is an equivalence relation on  $Q^n$

## Theorem

The set of first letter of an equivalence class is the whole stateset

$$u_1 u_2 \dots \equiv \mathbf{v} w u_1 \mathbf{v}' w' u_2 \dots$$

with  $\rho_{\mathbf{v}\mathbf{w}} = id$  and  $\rho_{\mathbf{v}'\mathbf{w}'} = id$

## Theorem [G., Klimann 2016]

A connected invertible-reversible Mealy automaton with a prime number of states does not generate an infinite Burnside group.

## Theorem [G., Klimann 2016]

A connected invertible-reversible Mealy automaton with a prime number of states does not generate an infinite Burnside group.

- ▶ Extend the result to any connected automaton
- ▶ Find an approach for non connected automata
- ▶ Decide the finiteness for automaton groups
- ▶ Find other structural properties

# Thanks