

Towards a better Understanding of the Scott Domain

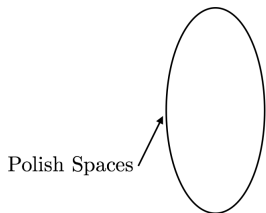
Louis Vuilleumier

joint work with Jacques Duparc

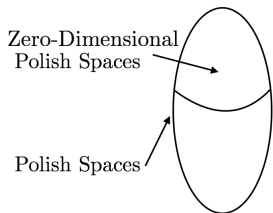
Department of Information Systems
University of Lausanne, Switzerland
Jacques.Duparc@unil.ch
Louis.Vuilleumier.1@unil.ch

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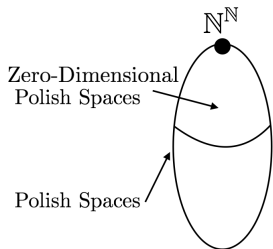
Classical Descriptive Set Theory



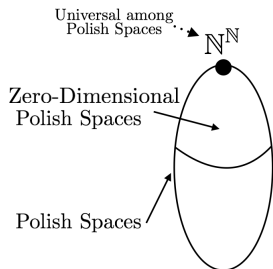
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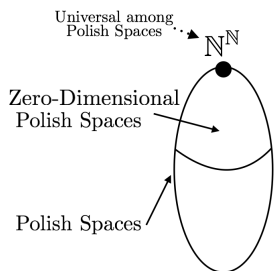
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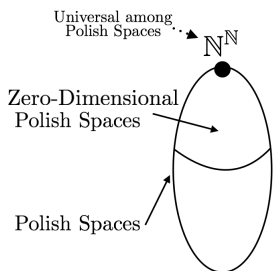


Classical Descriptive Set Theory



Study of the topological
complexity of the subsets

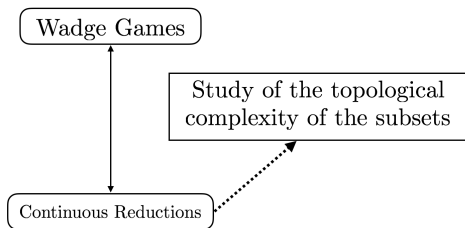
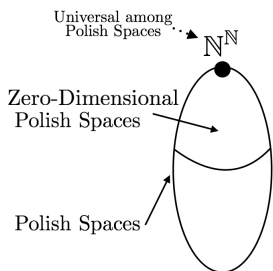
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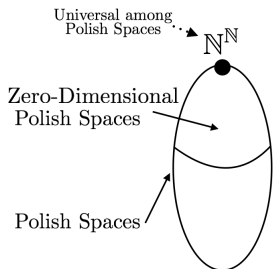
Continuous Reductions

Study of the topological
complexity of the subsets

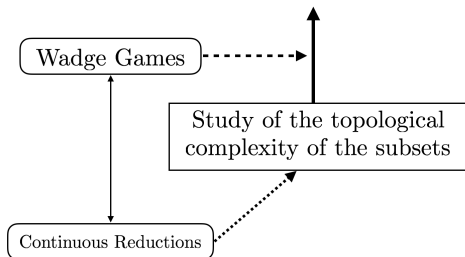
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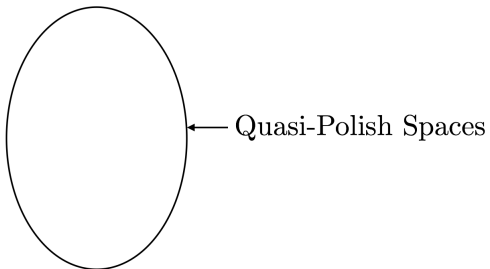
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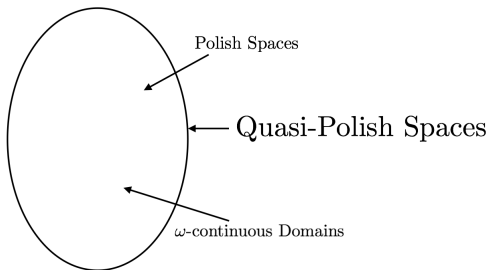
Full picture of the Wadge degrees of the Borel subsets of $\mathbb{N}^{\mathbb{N}}$



Generalization



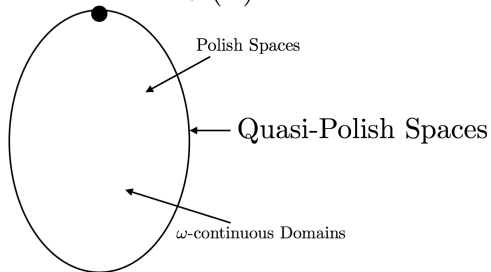
Generalization



Generalization

Universal among
Quasi-Polish Spaces

The Scott Domain $\mathcal{P}(\mathbb{N})$

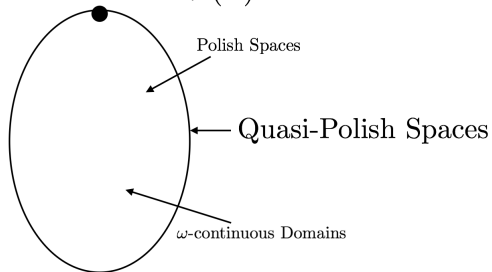


Generalization

Games on $\mathcal{P}(\mathbb{N})$?

Universal among
Quasi-Polish Spaces

The Scott Domain $\mathcal{P}(\mathbb{N})$



Scott domain and continuous reduction

Definition

Consider the set $\mathcal{P}(\mathbb{N})$ together with the topology — known as the **Scott topology** — generated by the basis

$$\mathcal{B} = \{\mathcal{O}_F : F \text{ finite subset of } \mathbb{N}\},$$

where

$$\mathcal{O}_F = \{X \subseteq \mathbb{N} : F \subseteq X\}.$$

This space is the **Scott domain**.

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A space is quasi-Polish if and only if it is homeomorphic to some $\mathcal{A} \in \Pi_2^0(\mathcal{P}(\mathbb{N}))$.

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Definition

Let $\mathcal{A}, \mathcal{B} \subseteq \mathcal{P}(\mathbb{N})$. If there exists a continuous function $f : \mathcal{P}(\mathbb{N}) \rightarrow \mathcal{P}(\mathbb{N})$ such that $f^{-1}(\mathcal{B}) = \mathcal{A}$, we say that \mathcal{A} is **continuously reducible** or **Wadge reducible** to \mathcal{B} , and we write $\mathcal{A} \leq_W \mathcal{B}$.

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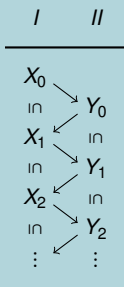
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\leq_W induces a quasi-order relation on the subsets of the Scott domain.

Games on the Scott domain

Definition

Let $\mathcal{A}, \mathcal{B} \subseteq \mathcal{P}(\mathbb{N})$. We define a game $G_\infty(\mathcal{A}, \mathcal{B})$.



$$X_n \in \mathcal{P}_{<\infty}(\mathbb{N}), Y_n \in \mathcal{P}(\mathbb{N})$$

II wins if and only if
 $(X \in \mathcal{A} \leftrightarrow Y \in \mathcal{B})$.

$$X = \bigcup_{n \in \mathbb{N}} X_n, Y = \bigcup_{n \in \mathbb{N}} Y_n$$

Strategy for \parallel

Definition

An **ultrapositional strategy** for \parallel in a game $G_\infty(\mathcal{A}, \mathcal{B})$ is an increasing function $\sigma : \mathcal{P}_{<\infty}(\mathbb{N}) \rightarrow \mathcal{P}(\mathbb{N})$, i.e. such that for all $X_0, X_1 \in \mathcal{P}_{<\infty}(\mathbb{N})$ with $X_0 \subseteq X_1$, we have $\sigma(X_0) \subseteq \sigma(X_1)$.

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An ultrapositional strategy is **winning** for II if, following this ultrapositional strategy and whatever I plays, II wins the game $G_\infty(\mathcal{A}, \mathcal{B})$.

Proposition

Let $\mathcal{A}, \mathcal{B} \subseteq \mathcal{P}(\mathbb{N})$. The following are equivalent.

- 1 $\mathcal{A} \leq_W \mathcal{B}$,
- 2 II has a winning ultrapositional strategy in $G_\infty(\mathcal{A}, \mathcal{B})$.

Examples

Consider $G_\infty(\{\mathbb{N}\}, \mathcal{P}_\infty(\mathbb{N}))$.

$$\begin{aligned} \sigma : \mathcal{P}_{<\infty}(\mathbb{N}) &\rightarrow \mathcal{P}(\mathbb{N}) \\ X &\mapsto \bigcup_{\substack{n \in \mathbb{N} \\ \{0, \dots, n\} \subseteq X}} \{0, \dots, n\}. \end{aligned}$$

It is a ultrapositional winning strategy.

Hence, $\{\mathbb{N}\} \leq_W \mathcal{P}_\infty(\mathbb{N})$.

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It is a ultrapositional winning strategy.

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Consider $G_\infty(\{\{0\}\}, \{\{0\}, \{0, 1, 2\}\})$.

$$\begin{aligned} \sigma : \mathcal{P}_{<\infty}(\mathbb{N}) &\rightarrow \mathcal{P}(\mathbb{N}) \\ \emptyset &\mapsto \emptyset, \\ \{0\} &\mapsto \{0\}, \\ X &\mapsto \{0, 1\} \quad \text{otherwise.} \end{aligned}$$

It is a ultrapositional winning strategy.

Hence, $\{\{0\}\} \leq_W \{\{0\}, \{0, 1, 2\}\}$.

Examples

Consider the game $G_\infty(\{\mathbb{N}\}, \{\{0\}\})$.

$$\begin{array}{cc} I & II \\ \hline \emptyset & \end{array}$$

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⋮

Hence, $\{\mathbb{N}\} \not\leq_w \{\{0\}\}$.

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⚡⚡⚡

Hence, $\{\mathbb{N}\} \not\leq_W \{\{0\}\}$.

We also have :

$(\{\mathbb{N}\}, \{\{0\}\}, \{\mathbb{N}\}^C, \{\{0\}\}^C)$ is an antichain with respect to the quasi-order \leq_W in the Scott domain.

No self-dual set in the Scott domain

Proposition

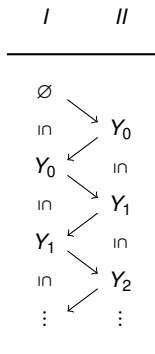
Let $\mathcal{A} \subseteq \mathcal{P}(\mathbb{N})$, then \mathcal{A} is non-self-dual.

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Proof.



$$X = \bigcup_{n \in \mathbb{N}} Y_n, \quad Y = \bigcup_{n \in \mathbb{N}} Y_n$$



Antichains in the Scott domain

Proposition (V.)

For every $k \in \mathbb{N} \setminus \{0\}$, there exists a sequence $(\mathcal{A}_1, \dots, \mathcal{A}_k)$ such that, for all $i, j \in \{1, \dots, k\}$, $i \neq j$, we have $\mathcal{A}_i \subseteq \Delta_3^0(\mathcal{P}(\mathbb{N}))$ and \mathcal{A}_i is not continuously reducible to \mathcal{A}_j .

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Merci !